

Use of Bits in 32 bit Word

Opava AP2, 06/04/2017

See:

www.cprogramming.com/tutorial/bitwise_operators.html

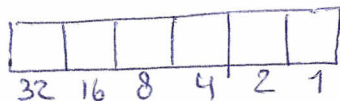
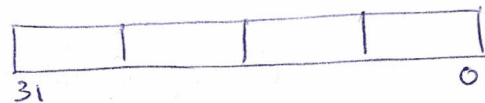
- ① What is int, short, long
- ② Operations on bits
- ③ Example: "Mini Dict".
- ④ Example: Blocks in Bloxor2.
- ⑤ Arithmetic by bits.

① Datatypes int, etc.

Manual: stores integer number between
-2147483647 & +2147483647

Memory (chips) store blocks of 32 bits. (four Bytes)

Representation: binary coding!



The "value" of the i^{th} bit
is 2^i , so eg 43 is

$$32 + 8 + 2 + 1 \text{ is } \boxed{101011}$$

With 6 bits used in this way, the max is $\boxed{111111}$,
so $32 + 16 + 8 + 4 + 2 + 1$, which is $\underline{2^6 - 1}$

Representation int uses bit nr 31 as sign (0 pos, 1 neg)
so 31 are used to store powers of two.

$$\Rightarrow \text{max is } \underline{2^{31} - 1} = 2147483648 \text{ or } \boxed{\text{Two Billion}}$$

①

② Programming language hides this from you!
Write all programs, compute all kinds of things
without knowing this!

Operators like $+$, $/$ are mapped to
processor instructions that treat the "word"
as an integer.

See www.cprogramming.com/tutorial/bitwise_operators.html
(says: space is usually not a problem
but I disagree!)

C, C++, C# and Java offer operations that
work directly on bits: extremely fast & powerful!

It works the same for short, int and long but my
examples are all just 8b long!

Shift Move bits to left or right,
fill on other side with zeroes.

$n =$

0	1	0	1	0	1	1	1
---	---	---	---	---	---	---	---

 (87)

$n \gg 3$

0	0	0	0	0	1	0	1	0
---	---	---	---	---	---	---	---	---

 (10)

The last three bits are gone!

So $(n \gg 3) \ll 3$ is not always the
same as n !

AND & takes two sequences, performs a pointwise AND

n

0	1	0	1	0	1	1
---	---	---	---	---	---	---

 (87)

m

1	1	1	1	0	0	0
---	---	---	---	---	---	---

 (240)

n & m

0	1	0	1	0	0	0
---	---	---	---	---	---	---

~~(240)~~ 80

87 & 240 = ~~240~~ 80

OR pointwise OR:

n 0 1 0 1 0 1 1 1

m 1 1 1 1 0 0 0 0

n | m 1 1 1 1 0 1 1 1 (243)

XOR Pointwise XOR: place 1 where there is exactly ONE 1

n 0 1 0 1 0 1 1 1

m 1 1 1 1 0 0 0 0

n ^ m 1 0 1 0 0 1 1 1 ^

NOT: unary operator! Bitwise complement

n

0	1	0	1
---	---	---	---

0	1	1	1
---	---	---	---

~n

1	0	1	0
---	---	---	---

1	0	0	0
---	---	---	---

③ Subset dictionary

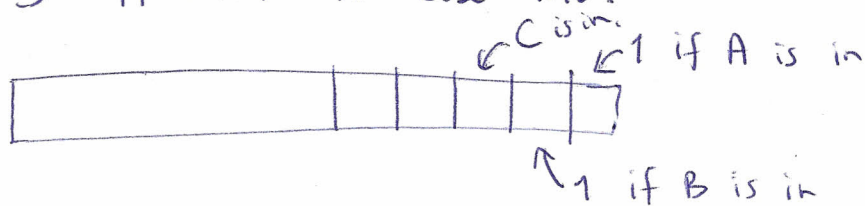
Program reasons about 20 objects A B C ...
and must store subsets. like { A D E F } set1

You can use a map
or a string, but it

{ A B } set2

{ } (= $\emptyset$) emp.

is very efficient to use int!



So set1 = $1 + 8 + 16 + 32 = 57$

6	5	4	3	2	1	0
0	0	1	1	1	0	1
G	F	E	D	C	B	A

set2 = 11 = 3

emp 000000 = 0.

3.1 See if object i is in: $((\text{set} \gg i) \& 1 == 1)$

e.g. set1 contains D, object nr 3!

set1 is 00111001

set1 >> 3 is 00000111

1 is 00000001

AND 00000001

which is indeed == 1

set2 does not contain D.

3.2 Insert item i in a set:

$$\text{set2} = \text{set} \mid (1 \ll i)$$

eg $\text{set1} \cup \{C\}$

You "want" $\text{set1}[2] = \text{True}$

$$\begin{array}{rcl} 1 & \text{is} & 00000001 \\ 1 \ll 2 & \text{is} & 00000100 \\ \text{set1} & \text{is} & 00111001 \\ \hline \text{OR of these two} & & 00111101, \end{array}$$

so looks "the same as" set1 but now with bit2 = 1.

3.3 Delete item i:

$$\text{set3} = \text{set2} \& \sim(1 \ll i)$$

How does this work?

$$\begin{array}{rcl} 1 \ll i & \text{is} & \boxed{00001000} \\ & & \quad \quad \quad i \\ \sim(1 \ll i) & \text{is} & \boxed{11110111} \end{array}$$

If I & this with some bitrow, all bits will remain except bit i which will be forced to 0.

$$\text{set2} = \boxed{00111101} \quad \begin{array}{l} \nwarrow \& \\ \searrow \end{array}$$

$$\text{set3} = \boxed{00110101}$$

3.4 ~~check~~ Compute $a \cup b : a \mid b$

3.5 Check if $a \subseteq b$ $(a \& \sim b) == 0$

} Think of examples yourself!

④ Packing values: Blocks of bits.

In Bloxor2, we have three parameters (i, j, or , four if we count the switches!) that are each small.

Assume i and j are ~~$\leq$~~ < 8 (3 bits)
or < 4 (2 bits).

Typical example

$i = 5$	$j = 7$	$or = 2$
---------	---------	----------

As bits: $i =$

0	0	0	0	0	1	0	1
---	---	---	---	---	---	---	---

Because we KNOW that $i < 8$, these bits are ALWAYS 0 so it isn't necessary to store them!

$j =$

0	0	0	0	0	1	1	1
---	---	---	---	---	---	---	---

↑ Five irrelevant bits of j

$or =$

0	0	0	0	0	0	1	0
---	---	---	---	---	---	---	---

Storing the three values as three numbers wastes a lot of space! Use an int (32b in my examples) to store the relevant bits only: $3+3+2$ is 8 so this fits!

$p =$

i	j	or
101	111	10

$$(190 = 32 \times 5 + 4 \times 7 + 2)$$

How can we form this number quickly?

I can left-shift i over 5 positions without losing information because I know the left 5 bits are 0:

$i \ll 5$ is

1	0	1	0	0	0	0	0
---	---	---	---	---	---	---	---

Similar

$j \ll 2$ is

0	0	0	1	1	1	0	0
---	---	---	---	---	---	---	---

and

or

is

0	0	0	0	0	0	1	0
---	---	---	---	---	---	---	---

160

28

2

+

190.

I can OR these strings:

$p =$

1	0	1	1	1	1	0	0
---	---	---	---	---	---	---	---

Addition works differently from OR, but as long as no column contains more 1's the outcome is the same!

So $p = (i \ll 5) | (j \ll 2) | \text{or}$

$= (((i \ll 3) | j) \ll 2) | \text{or}$

Find back j : $(p \gg 2) \& m3$

where $m3$ is a "mask" of three 1's:

0	0	0	0	0	1	1	1
---	---	---	---	---	---	---	---

which is the value $2^3 - 1$, or 7.

⑤ Arithmetic

Some simple arithmetic can be done with bit operations, especially when powers of 2 are involved!

Even $(x \% 2) == 0$ requires division ~~\$\$~~
\$

$(x \& 1) == 0$ only looks at last bit.

Multiply by 2 $2i$ is $2 * i$ (multiply)

or $i + i$ (add)

or $i \ll 1$ bitshift.

$2i+1$ is
 $(i \ll 1) | 1$

Multiply by 2^t : $2^t * a$ is $a \ll t$

Divide by 2^t $a / 2^t$ is $a \gg t$ (integer!)

Round down to multiple of 2^t eg with $t=3$,
 $85 \rightarrow 80$

$a \& \sim m3$ or $(a \gg t) \ll t$

⑥ See all bits:

for ($i = 31$; $i \geq 0$; $i--$)

if $((n \gg i) \& 1) == 1$ print '1'

else print '0'.